# CHANGE AND DELAY CONTRACTS FOR HYBRID SYSTEM COMPONENT VERIFICATION

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> Carnegie Mellon University

## **OVERVIEW**

Background

- □ Cyber-Physical Systems
- □ Component-based Verification
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  - □ Components
  - $\Box$  Contracts
  - $\Box$  Composition
- Evaluation
  - $\Box$  Implementation
  - $\Box$  Experiments

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- Model and analyze CPS: Hybrid system models
  - $\hfill\square$  Hybrid programs: program notation for CPS
  - $\Box$  Safety Property:  $\Phi \rightarrow [\alpha] \Psi$ 
    - System **contract**: Starting in  $\Phi$ , each run of hybrid program  $\alpha$  leads to a safe state  $\Psi$
    - Verified using hybrid systems theorem prover KeYmaera X
  - □ Analysis is challenging for **large monolithic models**

#### Idea: Component-based approach to hybrid system safety verification



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Monolithic

Internal Behavior System Component A || Component B )



Monolithic

External Behavior System Contract Contract A ∧ Contract B )

Internal Behavior System Component A || Component B )



Monolithic





Component-based

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- plant
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  else speed := 0
- $\approx$  pos'(t) = speed



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```
program \approx ( ctrl ; plant ; in ) *
```



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- □ Additional **assumptions** regarding initial state

#### Safety condition

□ Describes **safety** property and (optional) guarantees for values produced on ports

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contract  $\approx$  *initial*  $\rightarrow$  [ program ] *post* 

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    - Time passes simultaneously
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  - □ Compatible Communication
    - Robot expects position of obstacle close to previous position

Composite program

 $prog_3 \approx \left( (ctrl_1; ctrl_2 \cup ctrl_2; ctrl_1); (plant_1 \parallel plant_2); in_1; in_2; in_{open} \right) *$ 

Composite contract

 $cont_3 \approx (init_1 \wedge init_2) \rightarrow [prog_3](post_1 \wedge post_2)$ 

#### Theorem

- Composite program of two compatible components obeys composite contract
- $\hfill\square$  User provides  $\mbox{Proof A}$  and  $\mbox{Proof B}$
- □ Theorem derives **System Proof**
- → Safety verification results about contracts for components transfer to composites!

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- □ Contract, i.e., pre- and post condition
- Tactic to proof that components obey respective contracts

- Output
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## **EVALUATION – EXPERIMENTS**

Existing Case Studies
 Robot Collision Avoidance – Robix
 European Train Control System – ETCS
 Adaptive Cruise Control – LLC

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	Non-	Manual steps		Duration [s]	
linea	linear	Comp	Mono	Comp	Mono
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Summary

- □ Reduction of verification time
  - (especially for automated
  - proofs)
- $\hfill\square$  Reduction of proof effort

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