On a Decidable Fragment of $d\mathcal{L}$ or, The Next 700 (Un)decidable Fragments of $d\mathcal{L}$

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December 11, 2018



If you or a loved one has been frustrated trying to formally verify systems,



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you may be entitled to righteous indignation.

Why is formal verification so frustrating?

- complicated and tedious proofs
- lots of work for no product change
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Why is formal verification so frustrating?

- complicated and tedious proofs
- lots of work for no user-facing change
- people only care it looks like it works Cyberphysical systems are life-critical!



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Results

- \bullet Found and implemented decidable fragments of d ${\cal L}$ to ease verifying cyberphysical systems
- Found undecidable/inter-decidable fragments of d $\mathcal L$ to ease future decidability research

(Un)decidability Results

Arithmetical Approaches

	Integer Arithmetic	$d\mathcal{L}$
positive ∃	MRDP's Diophantine	Post Correspondence
positive \forall	polynomial ID testing	extended Platzer-Tan
bounded	finitary checking	Post Correspondence
single variable	trivial	Post Correspondence
purely $+$	Presburger	Post Correspondence
purely $ imes$	Skolem	Post Correspondence

(Un)decidability Results



	d ${\cal L}$
without \cup	MRDP's Diophantine
without ;	piecewise constant derivative reachability
without *	(exponential) polynomial star-free
only :=	Post Correspondence
only $?(-)$	reduction to $FOL_\mathbb{R}$
only $x' = f(x) \& Q$	piecewise constant derivative reachability
simultaneously $[lpha] P \wedge \langle lpha angle P$	when $[\alpha]P$ is

How can this be used for theorem proving?

- Work with simple ODEs
- Human identifies loop invariant
- That's it! Everything else is free.

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Loop invariants?

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 - Loop invariants?
 - Encode integer arithmetic: undecidable
- Restrict to polynomial solutions of ODEs

Theorem (DAG condition) Given $S \equiv x'_i = e_1, \dots, x'_n = e_n$, let G be a digraph s.t. edge from $x'_i = e_i$ to $x'_j = e_j \iff x_i$ occurs in e_j Then, S has a polynomial solution \iff G is acyclic.

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Proof sketch.

Back-sub in the topological order of G.

 $\bullet\,\sim$ 500 lines in OCaml

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- \bullet Shallow embedding of d ${\cal L}$ using weak higher-order abstract syntax
- Polynomial manipulation and ODE solver
- Z3 for quantifier elimination

Polynomial Star-Free: Demo

Verifying
$$x \ge 0 \land v \ge 0 \land a \ge 0 \rightarrow [x' = v, v' = a] \ x \ge 0$$

Common.Valid "unsat\n((declare-fun _x0!0 () Real)\n(proof\n ((?x254 (* a _x0!0 _x0!0))\n (let ((?x257 (* (251 ?x257)))\n (let ((\$x287 (>= ?x260 0.0)))\n (

Conclusion and Future Work

• Survey of restrictions for (un)decidability

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- Survey of restrictions for (un)decidability
- Decision procedures for theorem proving



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