

07: Control Loops & Invariants

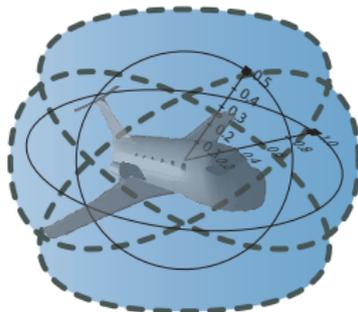
15-424: Foundations of Cyber-Physical Systems

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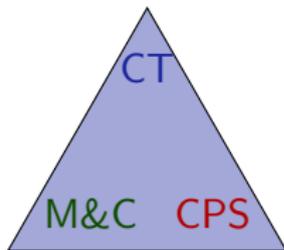
- 1 Learning Objectives
- 2 Induction for Loops
 - Iteration Axiom
 - Induction Axiom
 - Induction Rule for Loops
 - Loop Invariants
 - Simple Example
 - No Contexts for Soundness
- 3 Operationalize Invariant Construction
 - Bouncing Ball
 - Rescuing Misplaced Constants
 - Safe Quantum
- 4 Invariants
- 5 Summary

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Learning Objectives

Control Loops & Invariants

rigorous reasoning for repetitions
identifying and expressing invariants
global vs. local reasoning
relating iterations to invariants
finitely accessible infinities
operationalize invariant construction
splitting & generalizations



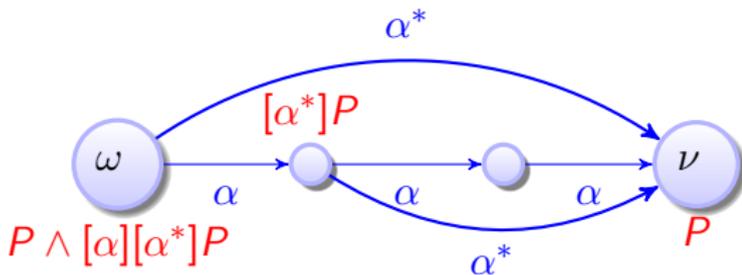
control loops
feedback mechanisms
dynamics of iteration

semantics of control loops
operational effects of control

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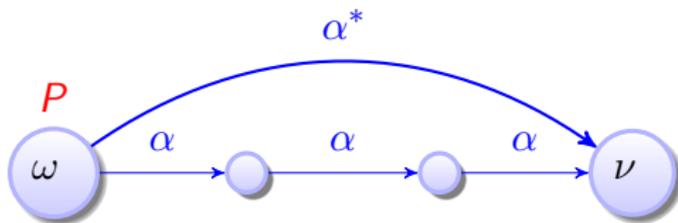
Iteration Axiom

$$[*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$



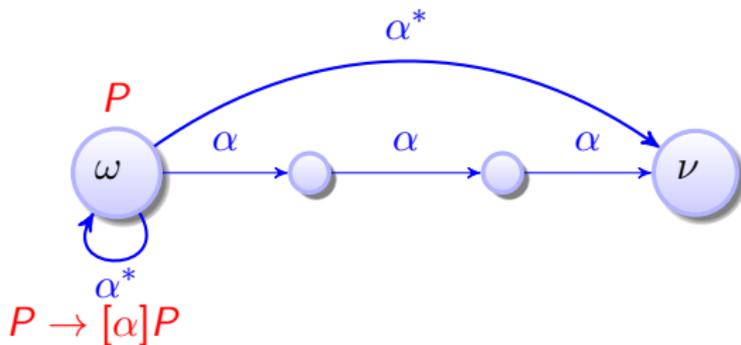
Lemma (I is sound)

$$I \ [\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$$



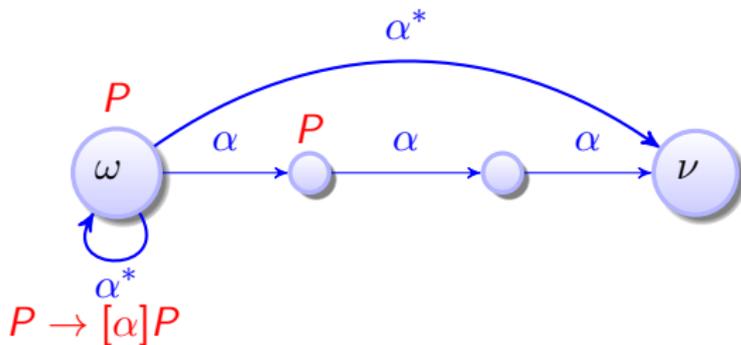
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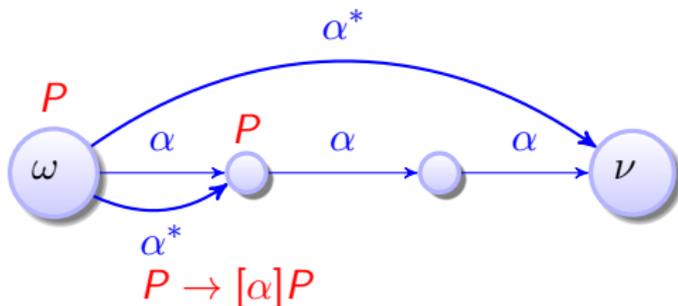
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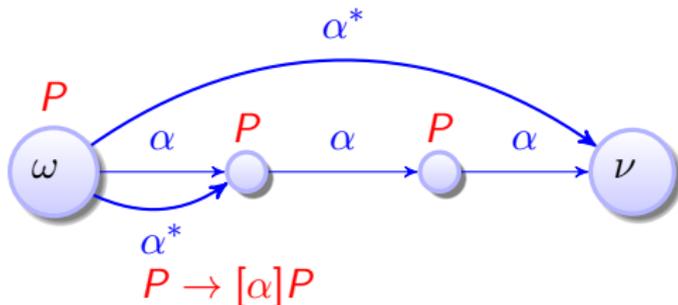
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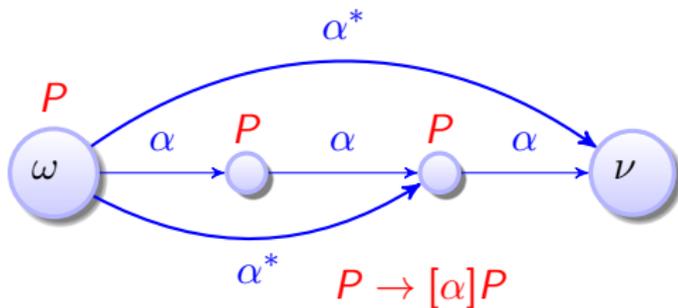
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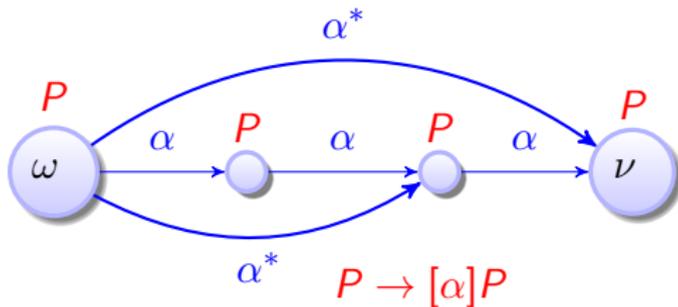
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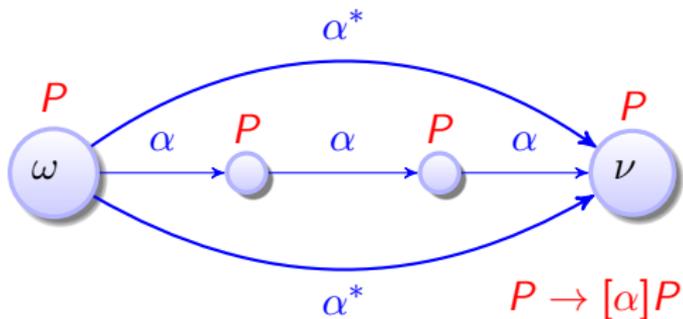
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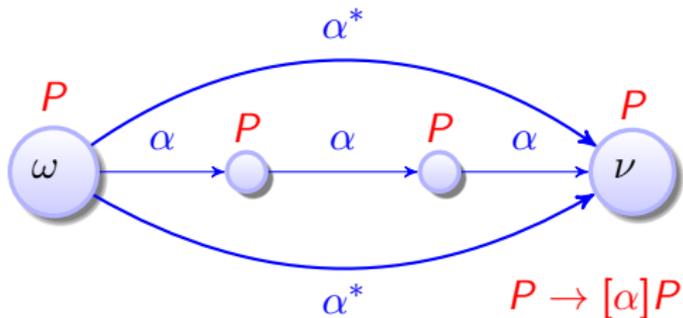
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Lemma (I is sound)

$$I \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$$



Problem: Inductive proof for $[\alpha^*]P$ needs proof of $[\alpha^*](P \rightarrow [\alpha]P)$

Induction Rule for Loops

Generalize induction step $[\alpha^*](P \rightarrow [\alpha]P)$ by Gödel

$$G \frac{P}{[\alpha]P}$$

Lemma (Loop induction rule *ind* is sound)

$$ind \frac{P \vdash [\alpha]P}{P \vdash [\alpha^*]P}$$

Induction Rule for Loops

Generalize induction step $[\alpha^*](P \rightarrow [\alpha]P)$ by Gödel

$$G \frac{P}{[\alpha]P}$$

Lemma (Loop induction rule *ind* is sound)

$$ind \frac{P \vdash [\alpha]P}{P \vdash [\alpha^*]P}$$

Proof (Derived rule).

$$\frac{\frac{id \frac{*}{P \vdash P} \quad \frac{\rightarrow R \frac{P \vdash [\alpha]P}{\vdash P \rightarrow [\alpha]P}}{G \frac{P \vdash [\alpha^*](P \rightarrow [\alpha]P)}}{\wedge R \frac{P \vdash P \wedge [\alpha^*](P \rightarrow [\alpha]P)}}}{I \frac{P \vdash [\alpha^*]P}}$$

□

Induction Rule for Loops

Generalize induction step $[\alpha^*](P \rightarrow [\alpha]P)$ by Gödel

$$G \frac{P}{[\alpha]P}$$

Lemma (Loop induction rule ind is sound)

$$ind \frac{P \vdash [\alpha]P}{P \vdash [\alpha^*]P}$$

Proof (Derived rule).

$$\frac{\frac{id \frac{*}{P \vdash P} \quad \frac{\rightarrow R \frac{P \vdash [\alpha]P}{\vdash P \rightarrow [\alpha]P}}{G \frac{P \vdash [\alpha^*](P \rightarrow [\alpha]P)}{P \vdash [\alpha^*](P \rightarrow [\alpha]P)}}{\wedge R \frac{P \vdash P \wedge [\alpha^*](P \rightarrow [\alpha]P)}}{I \frac{P \vdash [\alpha^*]P}}$$

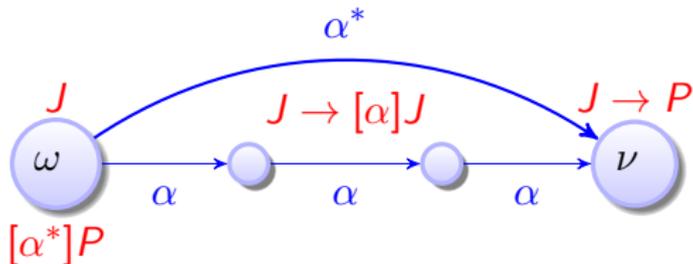
Problem: Rule ind is no equivalence. Its use of G may lose information:
 $[\alpha^*](P \rightarrow [\alpha]P)$ true but $P \vdash [\alpha]P$ is not valid. □

Loop Invariants

Generalize postcondition to strong loop invariant J by $M[\cdot]$ $\frac{P \rightarrow Q}{[\alpha]P \rightarrow [\alpha]Q}$

Lemma (Loop invariant rule loop is sound)

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$



Loop Invariants

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Proof (Derived rule).

$$\text{cut} \frac{\begin{array}{c} \text{ind} \frac{J \vdash [\alpha]J}{J \vdash [\alpha^*]J} \\ \rightarrow^R \frac{J \vdash [\alpha^*]J}{\Gamma \vdash J \rightarrow [\alpha^*]J, \Delta} \end{array} \quad \begin{array}{c} J \vdash P \\ \text{M}[\cdot] \frac{J \vdash P}{[\alpha^*]J \vdash [\alpha^*]P} \\ \rightarrow^L \frac{\Gamma \vdash J, \Delta \quad [\alpha^*]J \vdash [\alpha^*]P}{\Gamma, J \rightarrow [\alpha^*]J \vdash [\alpha^*]P, \Delta} \end{array}}{\Gamma \vdash [\alpha^*]P, \Delta}$$

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□

Problem: Finding invariant J can be a challenge.

Misplaced $[\alpha^*]$ suggests that J needs to carry along info about α^* history.

A Simple Discrete Loop Example

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\begin{array}{c} \text{loop} \\ \hline x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash J \quad J \vdash [x := x + y; y := x - 2 \cdot y]J \quad J \vdash x \geq 0 \\ \hline x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \vdash [(x := x + y; y := x - 2 \cdot y)^*] x \geq 0 \\ \hline \rightarrow R \\ \hline \vdash x \geq 8 \wedge 5 \geq y \wedge y \geq 0 \rightarrow [(x := x + y; y := x - 2 \cdot y)^*] x \geq 0 \end{array}$$

① $J \equiv x \geq 0$

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① $J \equiv x \geq 0$

stronger: Lacks info about y

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① $J \equiv x \geq 0$

stronger: Lacks info about y

② $J \equiv x \geq 8 \wedge 5 \geq y \wedge y \geq 0$

weaker: Changes immediately

A Simple Discrete Loop Example

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no: y may become negative if $x < y$

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no: y may become negative if $x < y$

④ $J \equiv x \geq y \wedge y \geq 0$

correct loop invariant

Add Context Γ, Δ to Premises?

$$\frac{\Gamma \vdash J, \Delta \quad \Gamma??, J \vdash [\alpha]J, \Delta?? \quad \Gamma??, J \vdash P, \Delta??}{\Gamma \vdash [\alpha^*]P, \Delta}$$

Add Context Γ, Δ to Premises?

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$$\begin{array}{l} \color{red}{\downarrow} \\ \frac{x = 0, x \leq 1 \vdash [x := x + 1]x \leq 1}{x = 0, x \leq 1 \vdash [(x := x + 1)^*]x \leq 1} \end{array}$$

Add Context Γ, Δ to Premises?

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Add Context Γ, Δ to Premises?

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Unsound! Be careful where your assumptions go.

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Proving Quantum the Acrophobic Bouncing Ball

$$A \vdash [(x'' = ; (?x=0; v := -cv \cup ?x \neq 0))^*] B(x, v)$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x, v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\text{loop} \frac{A \vdash j(x,v) \quad \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v) \quad j(x,v) \vdash B(x,v)}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

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Proving Quantum the Acrophobic Bouncing Ball

$$j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)$$

$$\text{loop} \frac{A \vdash j(x,v) \quad [:] \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash B(x,v)}}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

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Proving Quantum the Acrophobic Bouncing Ball

$$\frac{[:] \quad \frac{j(x,v) \vdash [x''=][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}}{[:]}$$

$$\frac{\text{loop} \quad A \vdash j(x,v) \quad [:] \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash B(x,v)}}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

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Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 j(x,v) \vdash [x''=]j(x,v) \\
 \hline
 j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v) \\
 \hline
 \text{MR} \frac{j(x,v) \vdash [x''=][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)} \\
 \text{[:] } \frac{}{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}
 \end{array}$$

$$\begin{array}{c}
 A \vdash j(x,v) \quad \text{[:] } \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)} \quad j(x,v) \vdash B(x,v) \\
 \hline
 \text{loop} \frac{}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}
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Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \text{MR} \\
 \text{[;]}
 \end{array}
 \frac{
 \frac{
 j(x,v) \vdash [x''=]j(x,v) \text{ [U]} \quad \frac{
 j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)
 }{
 j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v)
 }
 }{
 j(x,v) \vdash [x''=][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)
 }
 }{
 j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)
 }$$

$$\begin{array}{c}
 \text{loop}
 \end{array}
 \frac{
 A \vdash j(x,v) \text{ [;]} \quad \frac{
 j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v) \quad j(x,v) \vdash B(x,v)
 }{
 A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)
 }
 }{
 }$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x''=.. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \text{MR} \\
 \text{[;]} \\
 \hline
 \frac{j(x,v) \vdash [x''=]j(x,v) \quad \text{[U]} \quad \frac{\frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \quad j(x,v) \vdash [?x \neq 0]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)}}{\text{[;]} \quad \frac{j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}}{j(x,v) \vdash [x''=][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}}{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}
 \end{array}$$

$$\begin{array}{c}
 \text{loop} \\
 \hline
 \frac{A \vdash j(x,v) \quad \text{[;]} \quad \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v) \quad j(x,v) \vdash B(x,v)}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}
 \end{array}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \frac{j(x,v) \vdash [\text{?}x=0][v:=-cv]j(x,v)}{j(x,v) \vdash [\text{?}x=0; v:=-cv]j(x,v)} \text{[:]} \\
 \frac{j(x,v) \vdash [\text{?}x=0; v:=-cv]j(x,v) \quad j(x,v) \vdash [\text{?}x \neq 0]j(x,v)}{j(x,v) \vdash [\text{?}x=0; v:=-cv]j(x,v) \wedge [\text{?}x \neq 0]j(x,v)} \wedge R \\
 \frac{j(x,v) \vdash [x''=]j(x,v) \quad j(x,v) \vdash [\text{?}x=0; v:=-cv]j(x,v) \wedge [\text{?}x \neq 0]j(x,v)}{j(x,v) \vdash [x''=; (\text{?}x=0; v:=-cv \cup \text{?}x \neq 0)]j(x,v)} \text{[:]} \\
 \frac{j(x,v) \vdash [x''=][\text{?}x=0; v:=-cv \cup \text{?}x \neq 0]j(x,v)}{j(x,v) \vdash [x''=; (\text{?}x=0; v:=-cv \cup \text{?}x \neq 0)]j(x,v)} \text{[:]}
 \end{array}$$

$$\frac{A \vdash j(x,v) \quad j(x,v) \vdash B(x,v)}{A \vdash [(x''=; (\text{?}x=0; v:=-cv \cup \text{?}x \neq 0))^*]B(x,v)} \text{loop}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \frac{}{j(x,v), x=0 \vdash [v:=-cv]j(x,v)} \\
 \frac{[?]}{j(x,v) \vdash [?x=0][v:=-cv]j(x,v)} \\
 \frac{[!]}{j(x,v) \vdash [?x=0; v:=-cv]j(x,v)} \quad \frac{}{j(x,v) \vdash [?x \neq 0]j(x,v)} \\
 \wedge R \frac{}{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)} \\
 j(x,v) \vdash [x''=]j(x,v) \quad [U] \frac{}{j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v)} \\
 \hline
 MR \frac{}{j(x,v) \vdash [x''=][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)} \\
 [!]\frac{}{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}
 \end{array}$$

$$\text{loop} \frac{A \vdash j(x,v) \quad [!]\frac{}{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)} \quad j(x,v) \vdash B(x,v)}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \frac{j(x,v), x=0 \vdash j(x,-cv)}{[:=] \frac{j(x,v), x=0 \vdash [v:=-cv]j(x,v)}{[?] \frac{j(x,v) \vdash [?x=0][v:=-cv]j(x,v)}{[:] \frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v)}{\wedge R \frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}}} \\
 \frac{j(x,v) \vdash [x''=]j(x,v) \quad [U] \frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}}{MR \frac{j(x,v) \vdash [x''=][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}{[:] \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}}}
 \end{array}$$

$$\text{loop} \frac{A \vdash j(x,v) \quad [:] \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash B(x,v)}}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \frac{j(x,v), x=0 \vdash j(x,-cv)}{[:=] \frac{j(x,v), x=0 \vdash [v:=-cv]j(x,v)}{[?] \frac{j(x,v) \vdash [?x=0][v:=-cv]j(x,v)}{[:] \frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v)}{\wedge R \frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}}} \\
 \frac{j(x,v) \vdash [x''=]j(x,v)}{MR \frac{j(x,v) \vdash [x''=][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}{[:] \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}}}
 \end{array}$$

$$\text{loop} \frac{A \vdash j(x,v) \quad [:] \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash B(x,v)}}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \text{[:=]} \frac{j(x,v), x=0 \vdash j(x,-cv)}{j(x,v), x=0 \vdash [v:=-cv]j(x,v)} \\
 \text{[?]} \frac{j(x,v) \vdash [?x=0][v:=-cv]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv]j(x,v)} \quad \text{[?]} \frac{j(x,v), x \neq 0 \vdash j(x,v)}{j(x,v) \vdash [?x \neq 0]j(x,v)} \\
 \text{[?]} \frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv]j(x,v)} \quad \text{[?]} \frac{j(x,v), x \neq 0 \vdash j(x,v)}{j(x,v) \vdash [?x \neq 0]j(x,v)} \\
 \text{[?]} \frac{j(x,v) \vdash [?x=0; v:=-cv]j(x,v) \wedge [?x \neq 0]j(x,v)}{j(x,v) \vdash [?x=0; v:=-cv \cup ?x \neq 0]j(x,v)} \\
 \text{[?]} \frac{j(x,v) \vdash [x''=]j(x,v)}{j(x,v) \vdash [x''=][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)} \\
 \text{[?]} \frac{j(x,v) \vdash [x''=][?x=0; v:=-cv \cup ?x \neq 0]j(x,v)}{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}
 \end{array}$$

$$\text{loop} \frac{A \vdash j(x,v) \quad \text{[?]} \frac{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)}{j(x,v) \vdash [x''=; (?x=0; v:=-cv \cup ?x \neq 0)]j(x,v)} \quad j(x,v) \vdash B(x,v)}{A \vdash [(x''=; (?x=0; v:=-cv \cup ?x \neq 0))^*]B(x,v)}$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$A \vdash j(x, v)$$

$$j(x, v) \vdash [x'' = -g](j(x, v))$$

$$j(x, v), x = 0 \vdash j(x, -cv)$$

$$j(x, v), x \neq 0 \vdash j(x, v)$$

$$j(x, v) \vdash B(x, v)$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x, v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = -g \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash [\{x' = v, v' = -g \ \& \ x \geq 0\}](j(x,v))$$

$$j(x,v), x = 0 \vdash j(x, -cv)$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash [\{x' = v, v' = -g \ \& \ x \geq 0\}](j(x,v))$$

$$j(x,v), x = 0 \vdash j(x, -cv)$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

② $j(x,v) \equiv 0 \leq x \wedge x \leq H$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash [\{x' = v, v' = -g \ \& \ x \geq 0\}](j(x,v))$$

$$j(x,v), x=0 \vdash j(x,(-cv))$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

② $j(x,v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash \{ \{ x' = v, v' = -g \ \& \ x \geq 0 \} \} (j(x,v))$$

$$j(x,v), x=0 \vdash j(x,(-cv))$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

① $j(x,v) \equiv x \geq 0$

② $j(x,v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{ x' = v, v' = -g \ \& \ x \geq 0 \}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash [\{x' = v, v' = -g \ \& \ x \geq 0\}](j(x,v))$$

$$j(x,v), x=0 \vdash j(x,(-cv))$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

① $j(x,v) \equiv x \geq 0$

weaker: fails postcondition if $x > H$

② $j(x,v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash [\{x'=v, v'=-g \ \& \ x \geq 0\}](j(x,v))$$

$$j(x,v), x=0 \vdash j(x,(-cv))$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

① $j(x,v) \equiv x \geq 0$

weaker: fails postcondition if $x > H$

② $j(x,v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

③ $j(x,v) \equiv x = 0 \wedge v = 0$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash [\{x'=v, v'=-g \ \& \ x \geq 0\}](j(x,v))$$

$$j(x,v), x=0 \vdash j(x,(-cv))$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

① $j(x,v) \equiv x \geq 0$

weaker: fails postcondition if $x > H$

② $j(x,v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

③ $j(x,v) \equiv x = 0 \wedge v = 0$

strong: fails initial condition if $x > 0$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash [\{x'=v, v'=-g \ \& \ x \geq 0\}](j(x,v))$$

$$j(x,v), x=0 \vdash j(x,(-cv))$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

① $j(x,v) \equiv x \geq 0$

weaker: fails postcondition if $x > H$

② $j(x,v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

③ $j(x,v) \equiv x = 0 \wedge v = 0$

strong: fails initial condition if $x > 0$

④ $j(x,v) \equiv x = 0 \vee x = H \wedge v = 0$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x, v)$$

$$j(x, v) \vdash [\{x' = v, v' = -g \ \& \ x \geq 0\}](j(x, v))$$

$$j(x, v), x = 0 \vdash j(x, (-cv))$$

$$j(x, v), x \neq 0 \vdash j(x, v)$$

$$j(x, v) \vdash 0 \leq x \wedge x \leq H$$

- ① $j(x, v) \equiv x \geq 0$ weaker: fails postcondition if $x > H$
- ② $j(x, v) \equiv 0 \leq x \wedge x \leq H$ weak: fails ODE if $v \gg 0$
- ③ $j(x, v) \equiv x = 0 \wedge v = 0$ strong: fails initial condition if $x > 0$
- ④ $j(x, v) \equiv x = 0 \vee x = H \wedge v = 0$ no space for intermediate states

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x, v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash \{ \{x' = v, v' = -g \ \& \ x \geq 0\} \} (j(x,v))$$

$$j(x,v), x=0 \vdash j(x,(-cv))$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

① $j(x,v) \equiv x \geq 0$

weaker: fails postcondition if $x > H$

② $j(x,v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

③ $j(x,v) \equiv x = 0 \wedge v = 0$

strong: fails initial condition if $x > 0$

④ $j(x,v) \equiv x = 0 \vee x = H \wedge v = 0$

no space for intermediate states

⑤ $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash j(x,v)$$

$$j(x,v) \vdash \{ \{ x' = v, v' = -g \ \& \ x \geq 0 \} \} (j(x,v))$$

$$j(x,v), x=0 \vdash j(x,(-cv))$$

$$j(x,v), x \neq 0 \vdash j(x,v)$$

$$j(x,v) \vdash 0 \leq x \wedge x \leq H$$

① $j(x,v) \equiv x \geq 0$

weaker: fails postcondition if $x > H$

② $j(x,v) \equiv 0 \leq x \wedge x \leq H$

weak: fails ODE if $v \gg 0$

③ $j(x,v) \equiv x = 0 \wedge v = 0$

strong: fails initial condition if $x > 0$

④ $j(x,v) \equiv x = 0 \vee x = H \wedge v = 0$

no space for intermediate states

⑤ $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$

works: implicitly links v and x

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{ x' = v, v' = -g \ \& \ x \geq 0 \}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{aligned} &0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash 2gx = 2gH - v^2 \wedge x \geq 0 \\ &2gx = 2gH - v^2 \wedge x \geq 0 \vdash [\{x' = v, v' = -g \ \& \ x \geq 0\}](2gx = 2gH - v^2 \wedge x \geq 0) \\ &2gx = 2gH - v^2 \wedge x \geq 0, x = 0 \vdash 2gx = 2gH - (-cv)^2 \wedge x \geq 0 \\ &2gx = 2gH - v^2 \wedge x \geq 0, x \neq 0 \vdash 2gx = 2gH - v^2 \wedge x \geq 0 \\ &2gx = 2gH - v^2 \wedge x \geq 0 \vdash 0 \leq x \wedge x \leq H \end{aligned}$$

- ① $j(x,v) \equiv x \geq 0$ weaker: fails postcondition if $x > H$
- ② $j(x,v) \equiv 0 \leq x \wedge x \leq H$ weak: fails ODE if $v \gg 0$
- ③ $j(x,v) \equiv x = 0 \wedge v = 0$ strong: fails initial condition if $x > 0$
- ④ $j(x,v) \equiv x = 0 \vee x = H \wedge v = 0$ no space for intermediate states
- ⑤ $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$ works: implicitly links v and x

$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x,v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{aligned}0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 &\vdash 2gx = 2gH - v^2 \wedge x \geq 0 \\2gx = 2gH - v^2 \wedge x \geq 0 &\vdash [\{x' = v, v' = -g \ \& \ x \geq 0\}](2gx = 2gH - v^2 \wedge x \geq 0) \\2gx = 2gH - v^2 \wedge x \geq 0, x = 0 &\vdash 2gx = 2gH - (-cv)^2 \wedge x \geq 0 \\2gx = 2gH - v^2 \wedge x \geq 0, x \neq 0 &\vdash 2gx = 2gH - v^2 \wedge x \geq 0 \\2gx = 2gH - v^2 \wedge x \geq 0 &\vdash 0 \leq x \wedge x \leq H\end{aligned}$$

- ① $j(x, v) \equiv x \geq 0$ weaker: fails postcondition if $x > H$
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$$B(x, v) \equiv 0 \leq x \wedge x \leq H$$

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Proving Quantum the Acrophobic Bouncing Ball

$$\begin{aligned}0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 &\vdash 2gx = 2gH - v^2 \wedge x \geq 0 \\2gx = 2gH - v^2 \wedge x \geq 0 &\vdash [\{x' = v, v' = -g \ \& \ x \geq 0\}](2gx = 2gH - v^2 \wedge x \geq 0) \\2gx = 2gH - v^2 \wedge x \geq 0, x = 0 &\vdash 2gx = 2gH - (-cv)^2 \wedge x \geq 0 \quad \text{if } c = 1 \dots \\2gx = 2gH - v^2 \wedge x \geq 0, x \neq 0 &\vdash 2gx = 2gH - v^2 \wedge x \geq 0 \\2gx = 2gH - v^2 \wedge x \geq 0 &\vdash 0 \leq x \wedge x \leq H\end{aligned}$$

- ① $j(x, v) \equiv x \geq 0$ weaker: fails postcondition if $x > H$
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Proving Quantum the Acrophobic Bouncing Ball

$$\begin{aligned} &0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0 \vdash 2gx = 2gH - v^2 \wedge x \geq 0 \\ &2gx = 2gH - v^2 \wedge x \geq 0 \vdash [\{x' = v, v' = -g \ \& \ x \geq 0\}](2gx = 2gH - v^2 \wedge x \geq 0) \\ &2gx = 2gH - v^2 \wedge x \geq 0, x = 0 \vdash 2gx = 2gH - (-cv)^2 \wedge x \geq 0 \quad \text{if } c = 1 \dots \\ &2gx = 2gH - v^2 \wedge x \geq 0, x \neq 0 \vdash 2gx = 2gH - v^2 \wedge x \geq 0 \\ &2gx = 2gH - v^2 \wedge x \geq 0 \vdash 0 \leq x \wedge x \leq H \end{aligned}$$

- ① $j(x, v) \equiv x \geq 0$ weaker: fails postcondition if $x > H$
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$$A \equiv 0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge 1 \geq c \geq 0$$

$$B(x, v) \equiv 0 \leq x \wedge x \leq H$$

$$x'' = .. \equiv \{x' = v, v' = -g \ \& \ x \geq 0\}$$

Proving Quantum the Acrophobic Bouncing Ball

$$['] \quad j(x,v) \vdash [x'=v, v'=-g \ \& \ x \geq 0]j(x,v)$$

Proving Quantum the Acrophobic Bouncing Ball

$$\frac{[i] \quad j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt](x \geq 0 \rightarrow j(x,v))}{['] \quad j(x,v) \vdash [x' = v, v' = -g \ \& \ x \geq 0]j(x,v)}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{l} \text{[:=]} \\ \text{[:] } \\ \text{[']} \end{array} \frac{j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2][v := -gt](x \geq 0 \rightarrow j(x,v))}{j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt](x \geq 0 \rightarrow j(x,v))}$$
$$j(x,v) \vdash [x' = v, v' = -g \ \& \ x \geq 0]j(x,v)$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{l} \text{[:=]} \\ \hline j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2} t^2] (x \geq 0 \rightarrow j(x, -gt)) \\ \text{[:=]} \\ \hline j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2} t^2] [v := -gt] (x \geq 0 \rightarrow j(x,v)) \\ \text{[:]} \\ \hline j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2} t^2; v := -gt] (x \geq 0 \rightarrow j(x,v)) \\ \text{[:] } \\ \hline j(x,v) \vdash [x' = v, v' = -g \ \& \ x \geq 0] j(x,v) \end{array}$$

Proving Quantum the Acrophobic Bouncing Ball

$\forall R$	$j(x,v) \vdash \forall t \geq 0 (H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt))$
$[:=]$	$j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))$
$[:=]$	$j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2][v := -gt](x \geq 0 \rightarrow j(x,v))$
$[:]$	$j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt](x \geq 0 \rightarrow j(x,v))$
$[']$	$j(x,v) \vdash [x' = v, v' = -g \ \& \ x \geq 0]j(x,v)$

Proving Quantum the Acrophobic Bouncing Ball

$\rightarrow R$	$j(x,v) \vdash t \geq 0 \rightarrow H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt)$
$\forall R$	$j(x,v) \vdash \forall t \geq 0 (H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt))$
$[:=]$	$j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))$
$[:=]$	$j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2][v := -gt](x \geq 0 \rightarrow j(x,v))$
$[:]$	$j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt](x \geq 0 \rightarrow j(x,v))$
$[']$	$j(x,v) \vdash [x' = v, v' = -g \ \& \ x \geq 0]j(x,v)$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{l} j(x,v), t \geq 0, H - \frac{g}{2}t^2 \geq 0 \vdash j(H - \frac{g}{2}t^2, -gt) \\ \hline \rightarrow R \quad j(x,v) \vdash t \geq 0 \rightarrow H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt) \\ \hline \forall R \quad j(x,v) \vdash \forall t \geq 0 (H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt)) \\ \hline [:=] \quad j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2] (x \geq 0 \rightarrow j(x, -gt)) \\ \hline [:=] \quad j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2] [v := -gt] (x \geq 0 \rightarrow j(x,v)) \\ \hline [;] \quad j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt] (x \geq 0 \rightarrow j(x,v)) \\ \hline ['] \quad j(x,v) \vdash [x' = v, v' = -g \ \& \ x \geq 0] j(x,v) \end{array}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\overline{2gx=2gH-v^2 \wedge x \geq 0, H - \frac{g}{2}t^2 \geq 0 \vdash 2g(H - \frac{g}{2}t^2) = 2gH - (gt)^2 \wedge (H - \frac{g}{2}t^2) \geq 0}$$

$$\begin{array}{l} \overline{j(x,v), t \geq 0, H - \frac{g}{2}t^2 \geq 0 \vdash j(H - \frac{g}{2}t^2, -gt)} \\ \rightarrow R \quad \overline{j(x,v) \vdash t \geq 0 \rightarrow H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt)} \\ \forall R \quad \overline{j(x,v) \vdash \forall t \geq 0 (H - \frac{g}{2}t^2 \geq 0 \rightarrow j(H - \frac{g}{2}t^2, -gt))} \\ [:=] \quad \overline{j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2] (x \geq 0 \rightarrow j(x, -gt))} \\ [:=] \quad \overline{j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2] [v := -gt] (x \geq 0 \rightarrow j(x,v))} \\ [:] \quad \overline{j(x,v) \vdash \forall t \geq 0 [x := H - \frac{g}{2}t^2; v := -gt] (x \geq 0 \rightarrow j(x,v))} \\ ['] \quad \overline{j(x,v) \vdash [x' = v, v' = -g \ \& \ x \geq 0] j(x,v)} \end{array}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\frac{\overline{2gx=2gH-v^2} \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \quad \overline{H-\frac{g}{2}t^2 \geq 0} \vdash H-\frac{g}{2}t^2 \geq 0}{\wedge R \quad 2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0}$$

$$\frac{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)}{\rightarrow R \quad j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)}$$

$$\frac{\forall R \quad j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))}{[:=] \quad j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))}$$

$$\frac{[:=] \quad j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))}{[:] \quad j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))}$$

$$['] \quad j(x,v) \vdash [x'=v, v'=-g \ \& \ x \geq 0]j(x,v)$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \mathbb{R} \frac{\text{---} * \text{---}}{2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \quad H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0} \\
 \wedge \mathbb{R} \frac{\text{---}}{2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0} \\
 \frac{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)}{\rightarrow \mathbb{R} \frac{\text{---}}{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)}} \\
 \frac{\text{---}}{\forall \mathbb{R} \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))}} \\
 \frac{\text{---}}{[:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))}} \\
 \frac{\text{---}}{[:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))}} \\
 \frac{\text{---}}{[:] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))}} \\
 \frac{\text{---}}{['] \frac{\text{---}}{j(x,v) \vdash [x'=v, v'=-g \ \& \ x \geq 0]j(x,v)}}
 \end{array}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \mathbb{R} \frac{\text{---}^*}{2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2} \text{ id} \frac{\text{---}^*}{H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0} \\
 \wedge \mathbb{R} \frac{\text{---}}{2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0} \\
 \frac{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)}{\rightarrow \mathbb{R} \frac{\text{---}}{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)}} \\
 \frac{\text{---}}{\forall \mathbb{R} \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))}} \\
 \frac{\text{---}}{[:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))}} \\
 \frac{\text{---}}{[:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))}} \\
 \frac{\text{---}}{[i] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))}} \\
 \frac{\text{---}}{['] \frac{\text{---}}{j(x,v) \vdash [x'=v, v'=-g \ \& \ x \geq 0]j(x,v)}}
 \end{array}$$

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \mathbb{R} \frac{*}{2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2} \quad \text{id} \frac{*}{H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0} \\
 \wedge R \frac{2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0}{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)} \\
 \rightarrow R \frac{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)}{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)} \\
 \forall R \frac{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)}{j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))} \\
 [:=] \frac{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x, v))} \\
 [:=] \frac{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x, v))}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x, v))} \\
 ['] \frac{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x, v))}{j(x,v) \vdash [x'=v, v'=-g \ \& \ x \geq 0]j(x,v)}
 \end{array}$$

- Is Quantum done with his safety proof?

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \mathbb{R} \frac{\text{---}^*}{2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2} \quad \text{id} \frac{\text{---}^*}{H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0} \\
 \wedge \mathbb{R} \frac{\text{---}}{2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0} \\
 \frac{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)}{\rightarrow \mathbb{R} \frac{\text{---}}{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)}} \\
 \frac{\forall \mathbb{R} \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))}}{[:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))}} \\
 [:=] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))} \\
 [i] \frac{\text{---}}{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))} \\
 ['] \frac{\text{---}}{j(x,v) \vdash [x'=v, v'=-g \ \& \ x \geq 0]j(x,v)}
 \end{array}$$

- Is Quantum done with his safety proof?
- Oh no! The solutions we sneaked into ['] only solve the ODE/IVP if $x = 0, v = 0$ which $j(x,v)$ can't guarantee!

Proving Quantum the Acrophobic Bouncing Ball

$$\begin{array}{c}
 \frac{\mathbb{R} \frac{*}{2gx=2gH-v^2 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2} \quad \text{id} \frac{*}{H-\frac{g}{2}t^2 \geq 0 \vdash H-\frac{g}{2}t^2 \geq 0}}{\wedge R \frac{2gx=2gH-v^2 \wedge x \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash 2g(H-\frac{g}{2}t^2)=2gH-(gt)^2 \wedge (H-\frac{g}{2}t^2) \geq 0}}{j(x,v), t \geq 0, H-\frac{g}{2}t^2 \geq 0 \vdash j(H-\frac{g}{2}t^2, -gt)} \\
 \rightarrow R \frac{j(x,v) \vdash t \geq 0 \rightarrow H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt)}{\forall R \frac{j(x,v) \vdash \forall t \geq 0 (H-\frac{g}{2}t^2 \geq 0 \rightarrow j(H-\frac{g}{2}t^2, -gt))}{[:=] \frac{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2](x \geq 0 \rightarrow j(x, -gt))}{[:=] \frac{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2][v:=-gt](x \geq 0 \rightarrow j(x,v))}{[i] \frac{j(x,v) \vdash \forall t \geq 0 [x:=H-\frac{g}{2}t^2; v:=-gt](x \geq 0 \rightarrow j(x,v))}{['] \frac{j(x,v) \vdash [x'=v, v'=-g \ \& \ x \geq 0]j(x,v)}}
 \end{array}$$

- Is Quantum done with his safety proof?
- Oh no! The solutions we sneaked into ['] only solve the ODE/IVP if $x = 0, v = 0$ which $j(x,v)$ can't guarantee!
- **Never use solutions without proof!** \rightsquigarrow redo proof with true solution

loop $\frac{}{A \vdash [\alpha^*]B(x,v)}$

- 1 $j(x,v) \equiv 2gx=2gH-v^2 \wedge x \geq 0$
- 2 $p \equiv c=1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

$$\text{loop} \frac{\overline{\mathbb{R}} \quad A \vdash j(x,v) \wedge p \quad \square \wedge \quad \overline{j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p)}}{A \vdash [\alpha^*] B(x,v)}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

$$\text{loop} \frac{\mathbb{R} \frac{*}{A \vdash j(x,v) \wedge p} \quad \square \wedge \frac{}{j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p)}}{A \vdash [\alpha^*]B(x,v)}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Quantum Misplaced the Constants

$$\boxed{\wedge} [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\text{loop} \frac{\begin{array}{c} * \\ \hline \mathbb{R} A \vdash j(x,v) \wedge p \end{array} \quad \begin{array}{c} \wedge R \\ \hline j(x,v) \wedge p \vdash [\alpha]j(x,v) \wedge [\alpha]p \\ \hline \boxed{\wedge} \\ \hline j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p) \end{array}}{A \vdash [\alpha^*]B(x,v)}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Quantum Misplaced the Constants

$$\Box \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\frac{\text{loop} \quad \frac{\mathbb{R} \quad *}{A \vdash j(x,v) \wedge p} \quad \frac{\frac{\wedge R \quad \frac{\overline{j(x,v) \wedge p \vdash [\alpha]j(x,v)} \quad \overline{j(x,v) \wedge p \vdash [\alpha]p}}{j(x,v) \wedge p \vdash [\alpha]j(x,v) \wedge [\alpha]p}}{\Box \wedge \quad \frac{j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p)}{j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p)}} \quad \mathbb{R} \quad j(x,v) \wedge p \vdash B(x,v)}}{A \vdash [\alpha^*]B(x,v)}}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Quantum Misplaced the Constants

$$\Box \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\frac{\text{loop} \quad \frac{\mathbb{R} \quad *}{A \vdash j(x,v) \wedge p} \quad \frac{\text{above} \quad \frac{\frac{j(x,v) \wedge p \vdash [\alpha]j(x,v)}{\wedge R} \quad \frac{j(x,v) \wedge p \vdash [\alpha]p}{\vee}}{j(x,v) \wedge p \vdash [\alpha]j(x,v) \wedge [\alpha]p}}{\Box \wedge} \quad \mathbb{R} \quad j(x,v) \wedge p \vdash B(x,v)}{A \vdash [\alpha^*]B(x,v)}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Quantum Misplaced the Constants

$$\Box \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q \quad \forall p \rightarrow [\alpha]p \quad (FV(p) \cap BV(\alpha) = \emptyset)$$

$$\frac{\text{loop} \quad \frac{\mathbb{R} \quad *}{A \vdash j(x,v) \wedge p} \quad \frac{\text{above} \quad \frac{j(x,v) \wedge p \vdash [\alpha]j(x,v)}{\wedge R} \quad \frac{*}{j(x,v) \wedge p \vdash [\alpha]p}}{j(x,v) \wedge p \vdash [\alpha]j(x,v) \wedge [\alpha]p}}{\Box \wedge \quad \frac{j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p)}{\mathbb{R} \quad j(x,v) \wedge p \vdash B(x,v)}}}{A \vdash [\alpha^*]B(x,v)}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Quantum Misplaced the Constants

$$\Box \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q \quad \forall p \rightarrow [\alpha]p \quad (FV(p) \cap BV(\alpha) = \emptyset)$$

$$\frac{\text{loop} \quad \frac{\mathbb{R} \quad *}{A \vdash j(x,v) \wedge p} \quad \frac{\text{above} \quad \frac{j(x,v) \wedge p \vdash [\alpha]j(x,v)}{\wedge R} \quad \frac{*}{j(x,v) \wedge p \vdash [\alpha]p}}{j(x,v) \wedge p \vdash [\alpha]j(x,v) \wedge [\alpha]p} \quad \frac{*}{\mathbb{R} \quad j(x,v) \wedge p \vdash B(x,v)}}{\Box \wedge \quad \frac{j(x,v) \wedge p \vdash [\alpha](j(x,v) \wedge p)}{A \vdash [\alpha^*]B(x,v)}}$$

- 1 $j(x,v) \equiv 2gx = 2gH - v^2 \wedge x \geq 0$
- 2 $p \equiv c = 1 \wedge g > 0$
- 3 $J \equiv j(x,v) \wedge p$ as loop invariant

Quantum the Provably Safe Bouncing Ball

Proposition (Quantum can bounce around safely)

$$0 \leq x \wedge x = H \wedge v = 0 \wedge g > 0 \wedge \mathbf{1} = \mathbf{c} \rightarrow$$

$$[(x' = v, v' = -g \ \& \ x \geq 0; (?x = 0; v := -cv \cup ?x \neq 0))^*](0 \leq x \wedge x \leq H)$$

$$\text{@requires}(0 \leq x \wedge x = H \wedge v = 0)$$

$$\text{@requires}(g > 0 \wedge c = 1)$$

$$\text{@ensures}(0 \leq x \wedge x \leq H)$$

$$\{\{x' = v, v' = -g \ \& \ x \geq 0\};$$

$$(?x = 0; v := -cv \cup ?x \neq 0)\}^* \text{@invariant}(2gx = 2gH - v^2 \wedge x \geq 0)$$

Invariant Contracts

Invariants play a crucial role in CPS design. Capture them if you can. Use `@invariant` contracts in your hybrid programs.

Note: constants $c = 1 \wedge g > 0$ that never change are usually elided

- 1 Learning Objectives
- 2 Induction for Loops
 - Iteration Axiom
 - Induction Axiom
 - Induction Rule for Loops
 - Loop Invariants
 - Simple Example
 - No Contexts for Soundness
- 3 Operationalize Invariant Construction
 - Bouncing Ball
 - Rescuing Misplaced Constants
 - Safe Quantum
- 4 **Invariants**
- 5 Summary

The lion's share of understanding comes from understanding what does change (variants/progress measures) and what doesn't change (invariants).

Invariants are a fundamental force of CS

Variants are another fundamental force of CS

- 1 Learning Objectives
- 2 Induction for Loops
 - Iteration Axiom
 - Induction Axiom
 - Induction Rule for Loops
 - Loop Invariants
 - Simple Example
 - No Contexts for Soundness
- 3 Operationalize Invariant Construction
 - Bouncing Ball
 - Rescuing Misplaced Constants
 - Safe Quantum
- 4 Invariants
- 5 Summary

Summary

$$\vdash [\alpha^*]P \leftrightarrow P \wedge [\alpha^*](P \rightarrow [\alpha]P)$$

$$\text{loop} \frac{\Gamma \vdash J, \Delta \quad J \vdash [\alpha]J \quad J \vdash P}{\Gamma \vdash [\alpha^*]P, \Delta}$$

$$\text{G} \frac{P}{[\alpha]P}$$

$$\text{M}[\cdot] \frac{P \rightarrow Q}{[\alpha]P \rightarrow [\alpha]Q}$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\Box \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\forall p \rightarrow [\alpha]p \quad (FV(p) \cap BV(\alpha) = \emptyset)$$

- 6 Appendix
 - Iteration Axiom
 - Iterations & Splitting the Box
 - Iteration & Generalizations

compositional semantics \Rightarrow compositional rules!

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$A \vdash [\alpha^*]B$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\frac{\frac{[*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B}}{[*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}}{A \vdash [\alpha^*]B}}$$

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\begin{array}{c}
 \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)) \\
 \hline
 [*] \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 [*] \hline
 A \vdash B \wedge [\alpha][\alpha^*]B \\
 \hline
 [*] \hline
 A \vdash [\alpha^*]B
 \end{array}$$

Loops of Proofs: Iterations & Splitting the Box

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[\Box] \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{c}
 \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{[\Box] \wedge} \\
 \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \frac{}{[*]} \\
 \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{[*]} \\
 \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \frac{}{[*]} \\
 A \vdash [\alpha^*]B
 \end{array}$$

Loops of Proofs: Iterations & Splitting the Box

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[\wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{c}
 \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha]([\alpha]B \wedge [\alpha][\alpha][\alpha^*]B)} \\
 \frac{}{[\wedge} \\
 \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{[\wedge} \\
 \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \frac{}{[*]} \\
 \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{[*]} \\
 \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \frac{}{[*]} \\
 A \vdash [\alpha^*]B
 \end{array}$$

Loops of Proofs: Iterations & Splitting the Box

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[\wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{l}
 \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha]B \wedge [\alpha][\alpha][\alpha][\alpha^*]B} \\
 \frac{}{[\wedge} \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha]([\alpha]B \wedge [\alpha][\alpha][\alpha^*]B)} \\
 \frac{}{[\wedge} \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{[\wedge} \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \frac{}{[*]} \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \frac{}{[*]} \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \frac{}{[*]} \frac{}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Iterations & Splitting the Box

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[\Box] \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{l}
 \wedge R \frac{A \vdash B \quad A \vdash [\alpha]B \quad A \vdash [\alpha][\alpha]B \quad A \vdash [\alpha][\alpha][\alpha][\alpha^*]B}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha]B \wedge [\alpha][\alpha][\alpha][\alpha^*]B} \\
 [\Box] \wedge \frac{A \vdash B \wedge [\alpha]B \wedge [\alpha]([\alpha]B \wedge [\alpha][\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [\Box] \wedge \frac{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 [*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 [*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Iterations & Splitting the Box

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$[\Box] \wedge [\alpha](P \wedge Q) \leftrightarrow [\alpha]P \wedge [\alpha]Q$$

$$\begin{array}{c}
 A \vdash B \quad A \vdash [\alpha]B \quad A \vdash [\alpha][\alpha]B \quad A \vdash [\alpha][\alpha][\alpha][\alpha^*]B \\
 \wedge R \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha]B \wedge [\alpha][\alpha][\alpha][\alpha^*]B} \\
 [\Box] \wedge \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha]([\alpha]B \wedge [\alpha][\alpha][\alpha^*]B)} \\
 [\Box] \wedge \frac{}{A \vdash B \wedge [\alpha]B \wedge [\alpha][\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [\Box] \wedge \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 [*] \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [*] \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 [*] \frac{}{A \vdash [\alpha^*]B}
 \end{array}$$

- 1 Simple approach . . . if we don't mind unrolling until the end of time
- 2 Useful for finding counterexamples

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))$$

$$A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)$$

$$A \vdash B \wedge [\alpha][\alpha^*]B$$

$$A \vdash [\alpha^*]B$$

Loops of Proofs: Iterations & Generalizations

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c}
 A \vdash B \\
 \hline
 \wedge R \quad A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)) \\
 \hline
 [*] \quad A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)) \\
 \hline
 [*] \quad A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 [*] \quad A \vdash B \wedge [\alpha][\alpha^*]B \\
 \hline
 [*] \quad A \vdash [\alpha^*]B
 \end{array}$$

Loops of Proofs: Iterations & Generalizations

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c}
 A \vdash [\alpha]J_1 \quad \frac{}{J_1 \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 A \vdash B \text{ MR} \quad \frac{}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \wedge R \quad \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 [*] \quad \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [*] \quad \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 [*] \quad \frac{}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Iterations & Generalizations

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c}
 \begin{array}{c}
 A \vdash B \text{ MR} \frac{}{} \\
 \hline
 A \vdash [\alpha]J_1 \wedge R \frac{}{} \\
 \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 A \vdash B \wedge [\alpha][\alpha^*]B \\
 \hline
 A \vdash [\alpha^*]B
 \end{array} \\
 \hline
 \begin{array}{c}
 J_1 \vdash B \\
 \hline
 J_1 \vdash [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 J_1 \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)) \\
 \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)) \\
 \hline
 A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B) \\
 \hline
 A \vdash B \wedge [\alpha][\alpha^*]B \\
 \hline
 A \vdash [\alpha^*]B
 \end{array}
 \end{array}$$

Loops of Proofs: Iterations & Generalizations

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c}
 \begin{array}{c}
 J_1 \vdash [\alpha]J_2 \\
 \hline
 J_2 \vdash B \wedge [\alpha][\alpha^*]B
 \end{array} \\
 \text{MR} \frac{J_1 \vdash B \quad \text{---}}{J_1 \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \text{\color{green}AR} \frac{A \vdash [\alpha]J_1 \quad \text{---}}{J_1 \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \text{MR} \frac{A \vdash B \quad \text{---}}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \text{\color{green}AR} \frac{\text{---}}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 [*] \frac{\text{---}}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [*] \frac{\text{---}}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 [*] \frac{\text{---}}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Iterations & Generalizations

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$\begin{array}{c}
 J_2 \vdash B \quad \frac{}{J_2 \vdash [\alpha][\alpha^*]B} \\
 J_1 \vdash [\alpha]J_2 \wedge R \frac{}{J_2 \vdash B \wedge [\alpha][\alpha^*]B} \\
 J_1 \vdash B \text{MR} \frac{}{J_1 \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 A \vdash [\alpha]J_1 \wedge R \frac{}{J_1 \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 A \vdash B \text{MR} \frac{}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \wedge R \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 [*] \frac{}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 [*] \frac{}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 [*] \frac{}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Iterations & Generalizations

$$[*] [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$J_2 \vdash B \quad \frac{J_2 \vdash [\alpha]J_3 \quad \dots}{J_2 \vdash [\alpha][\alpha^*]B}$$

$$J_1 \vdash [\alpha]J_2 \quad \wedge\text{R} \frac{J_2 \vdash B \quad J_2 \vdash [\alpha][\alpha^*]B}{J_2 \vdash B \wedge [\alpha][\alpha^*]B}$$

$$J_1 \vdash B \quad \text{MR} \frac{J_1 \vdash [\alpha]J_2 \quad J_2 \vdash B \wedge [\alpha][\alpha^*]B}{J_1 \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$A \vdash [\alpha]J_1 \quad \wedge\text{R} \frac{J_1 \vdash B \quad J_1 \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}{J_1 \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$A \vdash B \quad \text{MR} \frac{A \vdash [\alpha]J_1 \quad J_1 \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$\wedge\text{R} \frac{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$[*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$[*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B}$$

$$[*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}$$

Loops of Proofs: Common Generalizations

$$\begin{array}{c}
 [*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P \\
 \\
 \text{MR} \quad \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta} \\
 \\
 J \vdash B \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B} \\
 \\
 J \vdash [\alpha]J \quad \wedge R \quad \frac{J \vdash B \quad \text{MR} \quad \frac{J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}}{J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \\
 A \vdash [\alpha]J \quad \wedge R \quad \frac{A \vdash B \quad \text{MR} \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \\
 \wedge R \quad \frac{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \\
 [*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \\
 [*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \\
 [*] \quad \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Extracting a Proof Rule

$$\begin{array}{c}
 \frac{J \vdash B}{A \vdash [\alpha^*]B} \\
 \\
 \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta} \text{MR} \\
 \\
 \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B} \\
 \\
 \frac{J \vdash [\alpha]J \quad \wedge R \frac{J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}}{J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)} \text{MR} \\
 \\
 \frac{A \vdash [\alpha]J \quad \wedge R \frac{J \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \text{MR} \\
 \\
 \wedge R \frac{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \\
 \\
 [*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \\
 \\
 [*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B} \\
 \\
 [*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}
 \end{array}$$

Loops of Proofs: Extracting a Proof Rule

$$\begin{array}{c}
 \frac{J \vdash [\alpha]J \quad J \vdash B}{A \vdash [\alpha^*]B} \quad \quad \quad [*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P \\
 \\
 \text{MR} \quad \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta} \\
 \\
 \begin{array}{c}
 J \vdash B \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B} \\
 \frac{J \vdash B \quad \text{MR} \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B}}{J \vdash B \wedge [\alpha][\alpha^*]B} \wedge R \\
 \frac{A \vdash [\alpha]J \quad \wedge R \quad \frac{J \vdash B \quad \text{MR} \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B}}{J \vdash B \wedge [\alpha][\alpha^*]B}}{A \vdash [\alpha]J \wedge [\alpha][\alpha^*]B} \wedge R \\
 \frac{A \vdash B \quad \text{MR} \quad \frac{J \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \wedge R \\
 \frac{\wedge R \quad \frac{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \wedge R}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))} \wedge R \\
 \frac{[*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \wedge R}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)} \wedge R \\
 \frac{[*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B} \wedge R}{A \vdash B \wedge [\alpha][\alpha^*]B} \wedge R \\
 \frac{[*] \quad \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B} \wedge R}{A \vdash [\alpha^*]B} \wedge R
 \end{array}
 \end{array}$$

Loops of Proofs: Extracting a Proof Rule

$$\frac{A \vdash J \quad J \vdash [\alpha]J \quad J \vdash B}{A \vdash [\alpha^*]B}$$

$$[*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \quad \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$J \vdash B \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B}$$

$$J \vdash [\alpha]J \quad \wedge\text{R} \quad \frac{J \vdash B \quad \text{MR} \quad \frac{J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash B \wedge [\alpha][\alpha^*]B}}{J \vdash B \wedge [\alpha][\alpha^*]B}$$

$$J \vdash B \quad \text{MR} \quad \frac{J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$A \vdash [\alpha]J \quad \wedge\text{R} \quad \frac{A \vdash [\alpha]J \quad J \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$A \vdash B \quad \text{MR} \quad \frac{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$\wedge\text{R} \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$[*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$[*] \quad \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B}$$

$$[*] \quad \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}$$

Loops of Proofs: Loop Invariants

$$\text{loop} \frac{A \vdash J \quad J \vdash [\alpha]J \quad J \vdash B}{A \vdash [\alpha^*]B}$$

Invariant J generalized
intermediate condition

$$[*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\text{MR} \quad \frac{\Gamma \vdash [\alpha]Q, \Delta \quad Q \vdash P}{\Gamma \vdash [\alpha]P, \Delta}$$

$$J \vdash B \quad \frac{J \vdash [\alpha]J \quad \dots}{J \vdash [\alpha][\alpha^*]B}$$

$$J \vdash [\alpha]J \quad \wedge R \frac{J \vdash [\alpha]J \quad J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash B \wedge [\alpha][\alpha^*]B}$$

$$J \vdash B \quad \text{MR} \frac{J \vdash B \quad J \vdash B \wedge [\alpha][\alpha^*]B}{J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$A \vdash [\alpha]J \quad \wedge R \frac{A \vdash [\alpha]J \quad J \vdash [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$A \vdash B \quad \text{MR} \frac{A \vdash B \quad A \vdash [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$\wedge R \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}$$

$$[*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B))}{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}$$

$$[*] \frac{A \vdash B \wedge [\alpha](B \wedge [\alpha][\alpha^*]B)}{A \vdash B \wedge [\alpha][\alpha^*]B}$$

$$[*] \frac{A \vdash B \wedge [\alpha][\alpha^*]B}{A \vdash [\alpha^*]B}$$



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