

Midterm I Exam

15-317/657 Constructive Logic
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February 13, 2020

Name: _____

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Instructions

- This exam is closed-book with one sheet of notes permitted.
- You have 80 minutes to complete the exam.
- There are 4 problems on 6 pages.
- Read each problem carefully before attempting to solve it.
- Do not spend too much time on any one problem.
- Consider if you might want to skip a problem on a first pass and return to it later.

	Max	Score
New Connections	45	
Harmonic Series	20	
Proof Terms	55	
VerifyThis Quantifier	30	
Total:	150	

1 New Connections (45 points)

Consider the new connective $\Box(A, B, C)$ that your friendly verificationists gave meaning to by the following introduction rule:

$$\frac{\frac{\overline{A \text{ true}} \quad u \quad \overline{B \text{ true}} \quad w}{\vdots} \quad \frac{\overline{B \text{ true}} \quad C \text{ true}}{\Box(A, B, C) \text{ true}}}{\Box I^{u,w}}$$

10 **Task 1** Give a set of elimination rules that harmoniously fit to $\Box I$:

10 **Task 2** Prove local soundness for the \Box connective.

10 **Task 3** Prove local completeness for the \Box connective.

5 **Task 4** Propose a proof term assignment for $\Box I^{u,w}$ **and** all other rules of the \Box connective.

$$\frac{\frac{\frac{}{ : A } \quad u \quad \frac{}{ : B } \quad w}{ : B } \quad \vdots}{ : B} \quad \frac{\frac{}{ : C } \quad w}{ : C} \quad \vdots}{ : \Box(A, B, C)} \Box I^{u,w}$$

10 **Task 5** Provide all local reduction rules for the proof terms of the \Box connective.

2 Harmonic Series (20 points)

This question considers introduction and elimination rules. Mark connectives as:

- Ⓜ for harmonious connectives and provide local reductions on proofs. You do not need to give local expansions (but convince yourself it is locally sound and locally complete).
- Ⓢ for unharmonious connectives and briefly explain one case that fails and why.

10 Task 1

$$\frac{A \text{ true} \quad B \text{ true}}{A \multimap B \text{ true}} \multimap I \quad \frac{A \multimap B \text{ true} \quad \frac{\frac{\overline{A \text{ true}}^u \quad \overline{B \text{ true}}^w}{\vdots} C \text{ true} \quad \frac{\overline{B \text{ true}}^w}{\vdots} C \text{ true}}{C \text{ true}} \multimap E^{u,w}}{C \text{ true}} \multimap E^{u,w}$$

10 Task 2

$$\frac{\frac{\overline{A \text{ true}}^u \quad \overline{B \text{ true}}^w}{\vdots} B \text{ true} \quad \frac{\overline{B \text{ true}}^w}{\vdots} A \text{ true}}{A \times B \text{ true}} \times I^{u,w} \quad \frac{A \times B \text{ true} \quad B \text{ true}}{A \text{ true}} \times E_1 \quad \frac{A \times B \text{ true} \quad A \text{ true}}{B \text{ true}} \times E_2$$

3 Proof Terms (55 points)

This question studies the proof terms of natural deduction. Recall that a proof term is called *normal*/irreducible if it cannot be reduced by any local reduction of proof terms.

10 **Task 1** Give a **normal** proof term justifying that $((A \supset B) \wedge (A \supset C)) \supset (A \supset (B \wedge C))$ is true.

10 **Task 2** What proposition is the following proof term justifying?
 $\text{fn } u \Rightarrow \text{fst}\langle \text{fn } v \Rightarrow \text{snd}(uv) \rangle, \text{ case inl } u \text{ of inl } w \Rightarrow (\text{fn } y \Rightarrow \text{fst}(uy)) \mid \text{inr } z \Rightarrow \text{abort } z \rangle$

10 **Task 3** Give a **normal** proof term justifying the proposition from Task 2 **or** explain why that proof term already is normal.

5 **Task 4** Give a **normal** proof term that justifies that $A \supset (A \vee B)$ is true.

10 **Task 5** Give a proof term that is **not** normal but justifies that $A \supset (A \vee B)$ is true.

10 **Task 6** Recall primitive recursion terms $R(n, t_0, x.r.t_s)$ for natural number n that are generated by successor s from 0. What **normal** proof term does the following proof term reduce to?
 $R(s\ 0, a, x.r.f(r, x, r))$

4 VerifyThis Quantifier (30 points)

Fill in missing **propositions** (if any), verification/use **judgments** (\uparrow and \downarrow) and **inference rule names** that make the following figures correct verifications **or** explain why that is **impossible** to do correctly. As in lecture, you do not need to decorate typing judgments $a : \tau$ with \uparrow or \downarrow .

15 Task 1

$$\begin{array}{c}
 \frac{}{\forall x:\tau. (A(x) \wedge B(x))} \quad u \quad \frac{}{a : \tau} \\
 \hline
 \frac{}{\quad} \\
 \hline
 \frac{A(a)}{\quad} \\
 \hline
 \forall x:\tau. A(x) \\
 \hline
 \frac{}{(\forall x:\tau. (A(x) \wedge B(x))) \supset (\forall x:\tau. A(x))}
 \end{array}$$

15 Task 2

$$\begin{array}{c}
 \frac{}{\exists x:\tau. C(x)} \quad u \quad \frac{\frac{}{C(a)} \quad w}{a : \tau C(a)}}{C(a)} \\
 \hline
 C(a) \\
 \hline
 (\exists x:\tau. C(x)) \supset C(a)
 \end{array}$$