

Part I

Proof Terms

Reduce the below proof terms taking care to not capture any substitutions!

- $$\begin{aligned}
 & (\lambda u : (A \supset B) \wedge (B \supset C). \lambda w : A. (\text{snd } u)((\text{fst } u)w))((\lambda x : A.x), \lambda y : B.y) \\
 & \Rightarrow_R \lambda w : A. (\text{snd } \langle (\lambda x : A.x), (\lambda y : B.y) \rangle)((\text{fst } \langle (\lambda x : A.x), (\lambda y : B.y) \rangle)w) \\
 & \Rightarrow_R \lambda w : A. (\lambda y : B.y)((\text{fst } \langle (\lambda x : A.x), (\lambda y : B.y) \rangle)w) \\
 & \Rightarrow_R \lambda w : A. (\lambda y : B.y)((\lambda x : A.x)w) \\
 & \Rightarrow_R \lambda w : A. (\lambda y : B.y)w \\
 & \Rightarrow_R \lambda w : A.w
 \end{aligned}$$

Note: This reduction showed that reducing the composition of two identity functions reduces again to the identity function.

- $$(\lambda y : A. \lambda f : A \supset B. fy)(fy).$$

We're renaming variables to avoid capture:

$$\begin{aligned}
 & (\lambda z : A. \lambda g : A \supset B. gz)(fy) \\
 & \Rightarrow_R \lambda g : A \supset B. g(fy)
 \end{aligned}$$

Part II

Tutch

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annotated proof andComm : A & B => B & A =
begin
[ u: A & B;
snd u: B;
fst u: A;
(snd u, fst u): B & A];
fn u => (snd u, fst u): A & B => B & A;
end;

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annotated proof mp: (A & (A => B)) => B =
begin
[ x: A & (A => B);
fst x: A;
snd x: A => B;
(snd x) (fst x): B];
fn x => (snd x) (fst x) : A & (A => B) => B;
end;

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annotated proof drop : A => (B => A) =
begin
[x: A;
[y: B;

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x: A];
fn y => x: B => A];
fn x => fn y=> x: A => (B => A);
end;

annotated proof impDef : (~A | B) => A => B =
begin
[ x: ~A | B;
[a: A;
[ na: ~A ;
na a: F;
abort (na a): B ];
[ b: B ;
b: B ];
case x of inl na => abort (na a) | inr b => b end: B ];
fn a => case x of inl na => abort (na a) | inr b => b end : A => B ];
fn x => fn a => case x of inl na => abort (na a) | inr b => b end : (~A | B) => A => B;
end;

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Part III

Verification

1. Give two different natural deduction proofs of $A \supset A \text{ true}$. How many natural deduction proofs of this judgement exist?

$$\frac{\overline{A \text{ true}}^u}{A \supset A \text{ true}} \supset I^u \qquad \frac{\frac{\overline{A \text{ true}}^u \quad \overline{A \text{ true}}^u}{A \wedge A \text{ true}} \wedge I \quad \frac{\overline{A \wedge A \text{ true}}}{A \text{ true}} \wedge E_L}{A \supset A \text{ true}} \supset I^u$$

There can be infinite proofs.

2. Give a proof of $A \supset A \uparrow$. How many proofs of this judgment exist?

$$\frac{\frac{\overline{A \downarrow}^u}{A \uparrow} \downarrow \uparrow}{A \supset A \uparrow} \supset I^u$$

Only one proof exists.